# Generic Double-Authentication Preventing Signatures and a Post-Quantum Instantiation

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Introduction

## **Digital Signatures**



Applications

- Signing transactions in cryptocurrencies
- Certificate and software signing
- $\cdot$  And many more

## Double-Authentication Preventing Signatures [PS14]



- Same context (address), different content
- » Extract secret key
  - Extraction from honest or malicious keys

Applications

- Deterring certificate subversion
- Double-spending prevention in offline payment channels
- Non-equivocation contracts

Approach	Address space	Extraction	Setting	Generic
[PS14]	exponential	DSE	factoring	×
[RKS15]	exponential	DSE	DLOG	×
[BPS17]	exponential	DSE	factoring	×
[BKN17]	exponential	DSE	SIS+LWE	×
[DRS18]	small	wDSE*	DLOG	$\checkmark$
[Poe18]	small	DSE	DLOG	×

Can we build efficient DAPS without address space limitation from existing signature schemes in a black-box way?

# DAPS without Structured Hardness Assumptions







 $\cdot$  One point reveals nothing about  $sk_{\Sigma}$ 



- Two points allow to recover  $\boldsymbol{s}\boldsymbol{k}_{\Sigma}$ 

 $\cdot\,$  DLOG-based signature scheme  $\Sigma$ 

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- » Polynomially sized address space

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- Signatures: secret share and consistency proof
- » Signature-of-knowledge style signature
- + Only requirement:  $\boldsymbol{\Sigma}$  public key image of one-way function









### commit to $\boldsymbol{\mathsf{sk}}_{\mathcal{F}}$





secret sharing of  $\boldsymbol{sk}_{\Sigma}$ 





#### consistency proof of PRF commitment



consistency proof of sharing polynomial computation



consistency proof of secret sharing

Generic approach:

- +  $\boldsymbol{\Sigma}$  with OWF relation between secret and public key
- + Verifiable Shamir secret sharing of  $\boldsymbol{\Sigma}$  secret key
- + Sharing polynomial determined by address

$$f_a(x) = \mathcal{F}(\mathsf{sk}_\mathcal{F}, a) \cdot x + \mathsf{sk}_\Sigma$$

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• Zero-knowledge proof of consistency

### Security (informal)

- Unforgeability from simulation-sound extractability, PRF and OWF properties
- Extraction from fixed-value-key-binding of PRF

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Extension

- + Extendable to N-authentication preventing signatures
- > Use degree N 1 sharing polynomial

### Instantiation

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- + OWF built from block cipher LowMC
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From stuctured hardness assumptions:

- + Fulfill secret-key-to-public-key relation requirement
- **?** Suitable proof system
- » Recent progress [AGM18]

## Extending Any Signature Scheme

# Signature scheme Σ Sign(Ҷ-Σ,...) Verify(⇔-Σ,...)

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- + From any other DAPS
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Security (informal)

- Unforgeability from unforgeability of signature scheme
- Extraction from extraction of DAPS

## Conclusion

Approach	Address space	Extraction	Setting	Generic
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[BPS17]	exponential	DSE	factoring	×
[BKN17]	exponential	DSE	lattices	×
[DRS18]	small	wDSE*	DLOG	$\checkmark$
[Poe18]	small	DSE	DLOG	×
Constr. 1	exponential	wDSE	symmetric	$\checkmark$
Constr. 2	exponential	DSE	any	$\checkmark$

#### Contribution

- ✓ Generic constructions of DAPS
- ✓ Construction 1: DAPS from symmetric-key primitives
- ✓ Construction 2: Extension of any signature scheme to DAPS

## Questions?

#### Full version: https://ia.cr/2018/790



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#### create offline payment channel

















receive deposit on misuse